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# Crypto 101: How to hide in plain sight

Batman's Kitchen 2024 | Dhruv Ashok

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## The basics: crypto vocab

- **Plaintext**: the message we want to communicate
- **Ciphertext**: the “scrambled” form of our plaintext/message
- **Key**: usually secret, used to convert plaintext to ciphertext and vice versa
- **Encryption**: the process of turning a plaintext into a ciphertext
- **Decryption**: the process of turning a ciphertext back into a plaintext
- **Nonce**: a (usually random) number only used once

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## The basics: encodings

- **Encoding != encryption**
- <https://cyberchef.io/>
- Hexadecimal to represent bytes, Base64 sometimes used as well
- Hex: **0x01020304050607**
- Base64: **dGVzdA==**

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## The basics: XOR

- Logical operation
- We like it because it's easily reversible
  - $ct = pt \text{ XOR key} \rightarrow pt = ct \text{ XOR key}$
- Used in a few different cryptosystems we'll talk about today
- Key properties
  - Anything XORed by itself is 0
    - Anything XOR 0 is itself
  - We can switch the order of stuff

Inputs		Output
A	B	X
0	0	0
0	1	1
1	0	1
1	1	0

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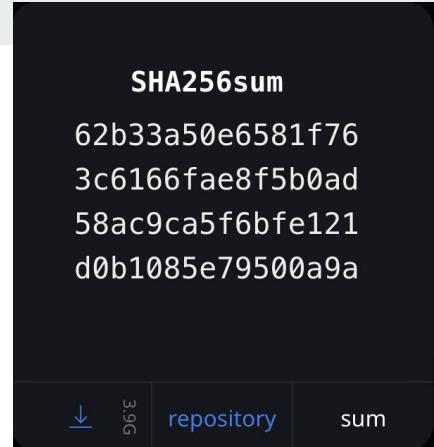
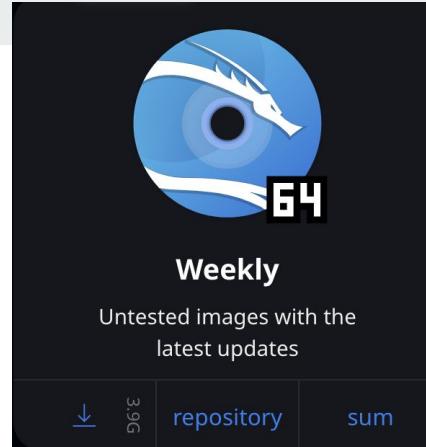
## The basics: hashes

- Take a variable length string, turn it into a fixed length string
- This is done **uniquely and irreversibly** (ideally)
- Few different algorithms for this
  - **Not all algorithms created equal!**
  - **MD5 bad, SHA1/SHA256 good**

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## The basics: hashes contd.

- File integrity
  - Provide hash of file, user can hash the file they've downloaded and check if it matches the legitimate hash
- Storing passwords
  - Hash and store a user's password when they create an account
  - For login, hash the password the user enters and check if it matches the stored one
  - Usually done in combination with a **salt**



## The basics: hashes (CTF version)

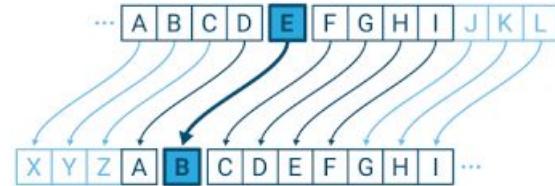
- We can't go from hash to plaintext..
  - But we can go from plaintext to hash
  - And with enough plaintext/hash pairs, we could get lucky!
- Hashcat, johntheripper, <https://crackstation.net>
- rockyou.txt

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# Encryption: from the beginning

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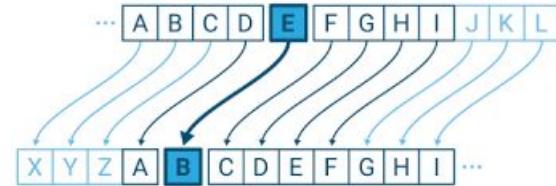
## Historical Ciphers: Caesar



- Choose a number between 1 and 26 (**key**)
- Given a message (**plaintext**), shift every alphabet to the right by the key
- Given a **ciphertext**, shift every alphabet to the left by the key
- $k = 3, m = \text{"HELLO"} \rightarrow ct = \text{"KHOOR"}$

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## Historical Ciphers: Caesar contd.

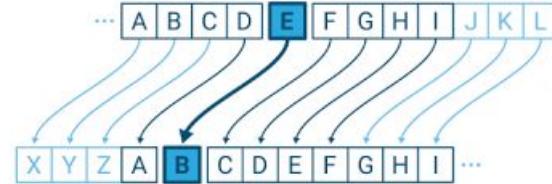


- <https://cyberchef.io>
- Choose ROT13 as the cipher
- Play around with it and get a feel for it
  - Given a ciphertext, can you determine anything about the plaintext **without** the key?

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## Caesar Cryptanalysis

- What are some problems with this?
  - Keyspace is small and easily brute forceable (1-26)
  - Frequency of characters in plaintext is leaked
  - Known plaintext issues



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# Historical Ciphers: Vigenere

- <https://cyberchef.io>
- Choose key (some string of alphabets)
- Repeat key until length of plaintext
- For each character pair, use table
- **Polyalphabetic** cipher

A	B	C	D	E	F	G	H	I	J	K	L	M	N	O	P	Q	R	S	T	U	V	W	X	Y	Z
A	B	C	D	E	F	G	H	I	J	K	L	M	N	O	P	Q	R	S	T	U	V	W	X	Y	Z
B	C	D	E	F	G	H	I	J	K	L	M	N	O	P	Q	R	S	T	U	V	W	X	Y	Z	A
C	D	E	F	G	H	I	J	K	L	M	N	O	P	Q	R	S	T	U	V	W	X	Y	Z	A	
D	E	F	G	H	I	J	K	L	M	N	O	P	Q	R	S	T	U	V	W	X	Y	Z	A	B	
E	F	G	H	I	J	K	L	M	N	O	P	Q	R	S	T	U	V	W	X	Y	Z	A	B	C	
F	G	H	I	J	K	L	M	N	O	P	Q	R	S	T	U	V	W	X	Y	Z	A	B	C	D	
G	H	I	J	K	L	M	N	O	P	Q	R	S	T	U	V	W	X	Y	Z	A	B	C	D	E	
H	I	J	K	L	M	N	O	P	Q	R	S	T	U	V	W	X	Y	Z	A	B	C	D	E	F	
I	J	K	L	M	N	O	P	Q	R	S	T	U	V	W	X	Y	Z	A	B	C	D	E	F	G	
J	K	L	M	N	O	P	Q	R	S	T	U	V	W	X	Y	Z	A	B	C	D	E	F	G	H	
K	L	M	N	O	P	Q	R	S	T	U	V	W	X	Y	Z	A	B	C	D	E	F	G	H	I	
L	M	N	O	P	Q	R	S	T	U	V	W	X	Y	Z	A	B	C	D	E	F	G	H	I	J	
M	N	O	P	Q	R	S	T	U	V	W	X	Y	Z	A	B	C	D	E	F	G	H	I	J	K	
N	O	P	Q	R	S	T	U	V	W	X	Y	Z	A	B	C	D	E	F	G	H	I	J	K	L	
O	P	Q	R	S	T	U	V	W	X	Y	Z	A	B	C	D	E	F	G	H	I	J	K	L	M	
P	Q	R	S	T	U	V	W	X	Y	Z	A	B	C	D	E	F	G	H	I	J	K	L	M	N	
Q	R	S	T	U	V	W	X	Y	Z	A	B	C	D	E	F	G	H	I	J	K	L	M	N	O	
R	S	T	U	V	W	X	Y	Z	A	B	C	D	E	F	G	H	I	J	K	L	M	N	O	P	
S	T	U	V	W	X	Y	Z	A	B	C	D	E	F	G	H	I	J	K	L	M	N	O	P	Q	
T	U	V	W	X	Y	Z	A	B	C	D	E	F	G	H	I	J	K	L	M	N	O	P	Q	R	
U	V	W	X	Y	Z	A	B	C	D	E	F	G	H	I	J	K	L	M	N	O	P	Q	R	S	
V	W	X	Y	Z	A	B	C	D	E	F	G	H	I	J	K	L	M	N	O	P	Q	R	S	T	
W	X	Y	Z	A	B	C	D	E	F	G	H	I	J	K	L	M	N	O	P	Q	R	S	T	U	
X	Y	Z	A	B	C	D	E	F	G	H	I	J	K	L	M	N	O	P	Q	R	S	T	U	V	
Y	Z	A	B	C	D	E	F	G	H	I	J	K	L	M	N	O	P	Q	R	S	T	U	V	W	
Z	A	B	C	D	E	F	G	H	I	J	K	L	M	N	O	P	Q	R	S	T	U	V	W	X	

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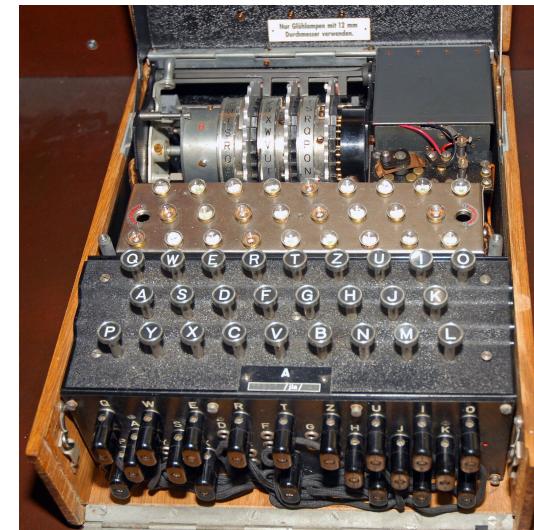
## Historical Ciphers: Vigenere contd.

- Keyspace is (if key chosen well) larger and harder to brute force
  - Still possible to do something like a dictionary attack
- Frequency is no longer leaked
- What happens if we want to scream though?
  - AAAAAAAAAAAAAAAA
- What happens if we know some of the plaintext?

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# Historical Ciphers: Enigma

- Employed by Nazis in WW2 to encrypt their communications
- Was famously broken after tireless cryptanalysis by Poland and Bletchley Park
- No letter **ever** encrypts to **itself**
  - Known plaintext was also taken advantage of



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## Historical Ciphers: One-Time Pad

- Choose key of length **at least** the length of the plaintext
- For each byte in plaintext and key, XOR them together
- pt = "encrypted", key = 0xb87301194dacf57fde → ct = 0xdd1d626b34dc811aba
- Looks pretty random when we have a good key!

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## Historical Ciphers: One-Time Pad contd.

- Employed by Soviets in WW2 to encrypt their communications
- OTP is actually quite good!....
  - But the Soviets misused it
  - Keys were repeated due to German invasion
- $(P1 \text{ XOR } k) \text{ XOR } (P2 \text{ XOR } k) = P1 \text{ XOR } P2$ 
  - **If you can get some/all of one, you can get some/all of the other since  $(P1 \text{ XOR } P2) \text{ XOR } P1 = P2$**
- This is actually how Soviet spies in Los Alamos were first discovered

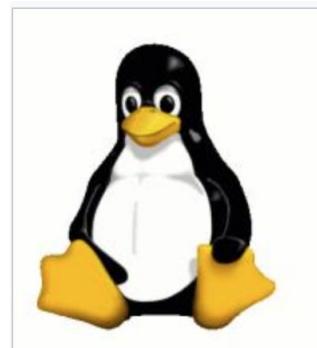
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# Symmetric key encryption

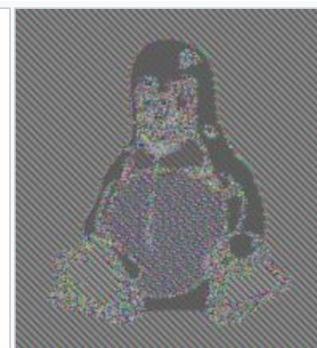
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## A better way: AES

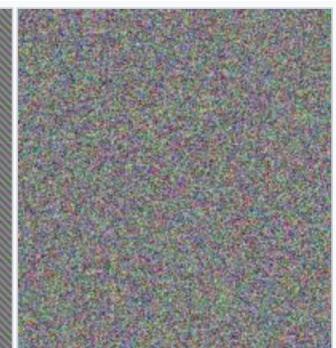
- Combines elements of the previous ciphers (OTP, substitution cipher, etc.)
- Has a few different **modes**
  - **Not all modes are created equal!**



Original image



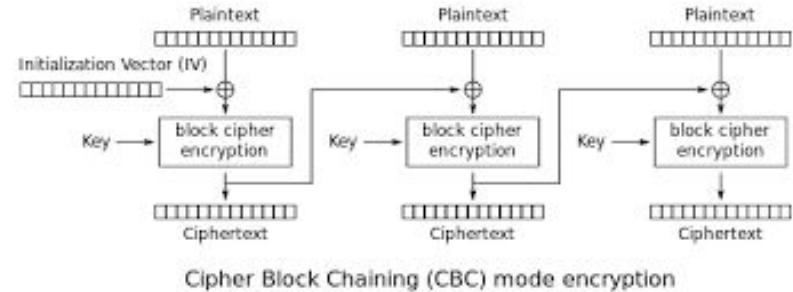
Encrypted using ECB mode



Modes other than ECB result in pseudo-randomness

## A better way: AES modes

- Block cipher (ECB, CBC, GCM)
  - Pad plaintext to a multiple of block size, then encrypt using notion of “blocks” (groups of bytes of the same size)
- Stream cipher (CTR)
  - Generate keystream of same size as plaintext, then XOR together
    - This is essentially the same as how OTP works!



## A better way: AES (CTF version)

- AES CTR key/IV reuse
  - Basically the same way you break One Time Pad
- AES CBC bitflipping
  - If you control the IV (which you shouldn't), you can predictably modify a decrypted ciphertext
- AES CBC padding oracle
  - Relies on excessive error details (padding failure)
- AES ECB chosen plaintext attack
- Various other shenanigans with how AES gets used/implemented

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## A better way: AES contd.

- We now introduce **randomness**
- Not vulnerable to frequency attacks, known plaintext (**usually**), key brute force (**usually**)
- It works beautifully but...
  - How do we communicate the key?
  - Especially if we're communicating for the first time on the Internet

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# Public key cryptography

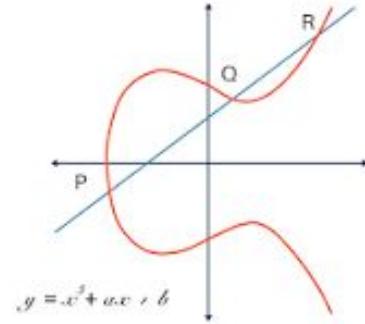


## RSA origin story

Ron Rivest, Adi Shamir, and Leonard Adleman at the Massachusetts Institute of Technology made several attempts over the course of a year to create a function that was hard to invert. Rivest and Shamir, as computer scientists, proposed many potential functions, while Adleman, as a mathematician, was responsible for finding their weaknesses. They tried many approaches, including "knapsack-based" and "permutation polynomials". For a time, they thought what they wanted to achieve was impossible due to contradictory requirements.<sup>[5]</sup> In April 1977, they spent Passover at the house of a student and drank a good deal of wine before returning to their homes at around midnight.<sup>[6]</sup> Rivest, unable to sleep, lay on the couch with a math textbook and started thinking about their one-way function. He spent the rest of the night formalizing his idea, and he had much of the paper ready by daybreak. The algorithm is now known as RSA – the initials of their surnames in same order as their paper.<sup>[7]</sup>

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## A new notion: public/private keys



- Symmetric crypto uses just one (secret) key to encrypt
- Asymmetric crypto uses a **public/private keypair**
  - You can share your public key with everyone!
  - NEVER share your private key with ANYONE.
- **Encrypt with public key, decrypt with private key**

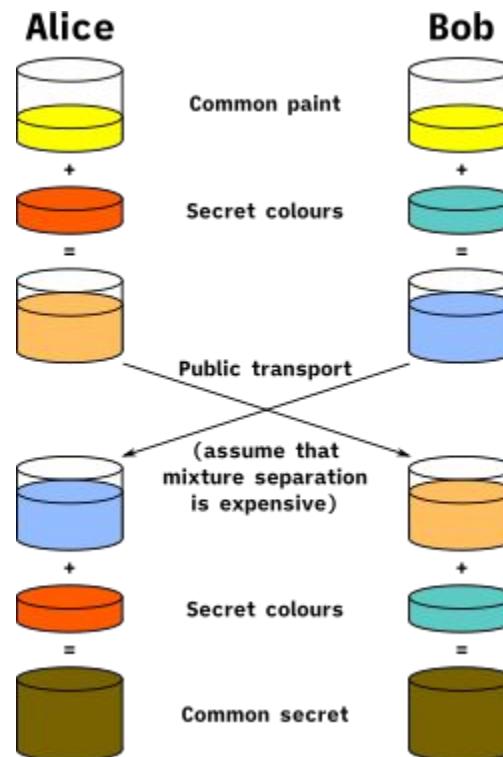
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## A new notion: public/private keys contd.

- We can use public key crypto to..
  - Encrypt communication
  - Sign/verify the authenticity of communication
    - This is how JSON Web Tokens (JWTs) work!
  - **Complete a key exchange**

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## Key exchange analogy



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## Why does this work?

- Public key cryptosystems are set up such that
  - encryption/decryption given the right information is easy
  - encryption/decryption is computationally infeasible otherwise
  - This idea is known as a “trapdoor function”
- Send information that helps us come to agreement on a key, but not enough that an attacker would be able to figure it out



## Public key crypto (CTF version)

- Main cryptosystem you will see is RSA (though many others exist, of course)
- <https://crypto.stanford.edu/~dabo/papers/RSA-survey.pdf>
- <https://github.com/RsaCtfTool/RsaCtfTool>
- Understanding the math of RSA will help a lot if you're interested in digging deeper
  - Feel free to ask me in more detail about this! :D

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# Further Resources

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## Challenges/resources

- <https://cryptohack.org>
- <https://cryptopals.com>
- A bunch of picoCTF challenges
  - Small point value ones cover some RSA and some ancient crypto
  - From there, some symmetric crypto/more complicated public key crypto/other random cool stuff
- CSE 426 (cryptography class here at UW, fairly theoretical)
- A lot of this content is covered in INFO 310 as well

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## Directions you can go with crypto

- As it relates to security
  - Ensure encryption is being used where appropriate and **being used properly**
- Theoretical
  - Constructing post-quantum cryptosystems
  - Constructing robust crypto that has specific properties
    - Multiparty computation (stuff that Dr. Benaloh talked about!)
- Practical
  - Side channel attacks
    - Power analysis
    - Timing
    - Heat
    - Electromagnetic radiation
    - ??? (dream it up, someone can probably figure out a way)

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Thank you!