

The Dark Arts of Hacking the “Unhackable”

Some random guy yapping about side channels and more useless junk you never knew you didn't not need

Batman's Kitchen
Feb 19th, 2025

Hacking is pretty simple in theory...

Some guy writes shitty code

You find a bug in it

???

Profit

```
#include <stdio.h>
#include <stdlib.h>
#include <string.h>

int main() {
    char true_pass[8] = "*****";
    char user_pass[8];
    gets(user_pass);

    int match = memcmp(user_pass, true_pass, 8);

    if (match != 0) {
        puts("incorrect password\n");
        exit(1);
    }

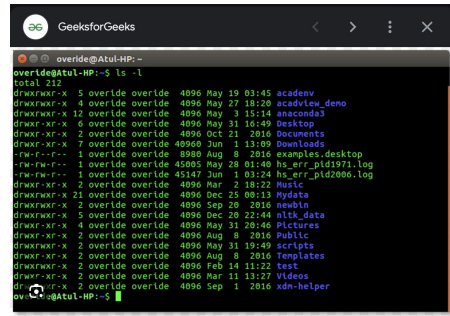
    puts("you win!");
    exec("/bin/sh");
}
```

Bugs

Never use **gets()**. Because it is impossible to tell without knowing the data in advance how many characters **gets()** will read, and because **gets()** will continue to store characters past the end of the buffer, it is extremely dangerous to use. It has been used to break computer security. Use **fgets()** instead.

^ from:

<https://linux.die.net/man/3/gets>



Hacking is pretty simple in theory...

Some guy writes shitty code

and a bug in it

???

P

```
#include <stdio.h>
#include <stdlib.h>
#include <string.h>

int main() {
    char true_pass[8] = "*****";
    char user_pass[8];
    fgets(user_pass, 8, stdin);

    int match = memcmp(user_pass, true_pass, 8);

    if (match != 0) {
        puts("incorrect password\n");
        exit(1);
    }

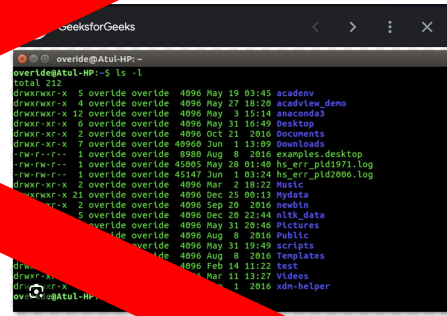
    puts("you win!");
    exec("/bin/sh");
}
```

Bugs

Never use **gets()**. Because it is impossible to tell without knowing the data in advance how many characters **gets()** will read, and because **gets()** will continue to store characters past the end of the buffer, it is extremely dangerous to use. It has been used to break computer security. Use **fgets()** instead.

^ from:

<http://man7.org/linux/man-pages/man3/fgets.3.html>





What is a hacker?

- Someone who exploits the gaps between assumptions and designs
- Not just computer systems!
 - Social systems (people and organizations)
 - Physical security systems
 - Public information systems

Challenging Computer Assumptions

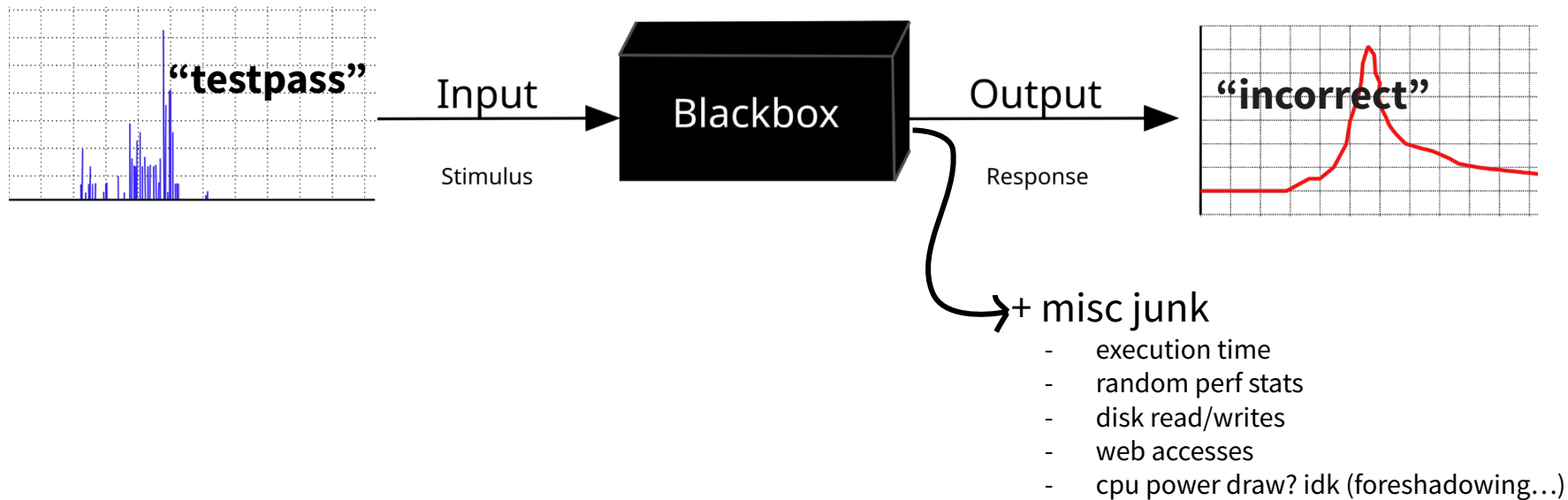
The language computers speak is a superset of the language we think computers speak



Challenging Computer Assumptions

The language computers speak is a superset of the language we think computers speak

??? ^ wtf does this mean



↻ You Retweeted



JavaScript Daily

@JavaScriptDaily



Presenting the "sleep sort", our favorite $O(\text{wtf})$ sorting algorithm. (C'mon, it's the weekend..) 😂

```
// sleep sort  
  
numbers = [8, 42, 38, 111, 2, 39, 1];  
  
numbers.forEach(num => {  
  setTimeout(() => { console.log(num) }, num);  
});
```

3:03 PM · Apr 23, 2017 · [Buffer](#)

1.2K Retweets 1.7K Likes



Woo! Let's do an activity

The screenshot shows a poker game interface with a dark blue background. On the left, a sidebar contains game information: the title "The Eye", a rule "No repeat hand types this round", a score requirement "Score at least 1,340,000 to earn \$\$\$\$\$", a "Round score" of 0, a "Push" button, a multiplier "340 X 21,600", a "Run Info" button, and an "Options" button. Below these are buttons for "Hands" (3), "Discards" (1), "Ante" (12/8), and "Round" (1). The main play area displays a "Joker" deck with six cards: a DNA helix, a blue joker, a purple joker, a green joker, a rainbow joker, and a white joker labeled "GAINO". Below the deck is a "X1.5 Mult" multiplier. The community cards are the Ace of Diamonds, King of Spades, Queen of Hearts, Queen of Clubs, and 10 of Diamonds. The player's hand consists of the 6 of Clubs, 3 of Spades, and 2 of Clubs. A discard pile on the right shows a patterned card. The interface includes progress indicators "6/7" and "0/3" for the deck and "3/8" and "5/52" for the player's hand.

The Eye
No repeat hand types this round
Score at least **1,340,000** to earn \$\$\$\$\$

Round score: **0**

Push
340 X 21,600

Run Info

Options

Hands: **3** Discards: **1**

\$4

Ante: **12/8** Round: **1**

JOKER 6/7

X1.5 Mult 0/3

Community cards: A♦, K♠, Q♥, Q♣, 10♦

Player's hand: 6♣, 3♠, 2♣ 3/8

Discard pile: 5/52

Case Study #1: Programs aren't instantaneous

poopoopeepee

“Secure” programs aren’t always safe

- The threat model programmers use != the real world threat model
- We’re slowly getting there! (... sort of)
- <https://www.chosenplaintext.ca/articles/beginners-guide-constant-time-cryptography.html>
- Are these attacks practical??



Breaking Bad: How Compilers Break Constant-Time Implementations

Moritz Schneider
ETH Zurich

Daniele Lain
ETH Zurich

Ivan Puddu
ETH Zurich

Nicolas Dutly
ETH Zurich

Srdjan Čapkun
ETH Zurich

Case Study #2: Web exploitation on secure pages

xslaeks

Web exploitation without javascript??

README.md



sanitizer - 500 points - 1 solve

What does it truly mean to sanitize something? What is a sanitize? Is this challenge solvable?

Admin bot code is provided, running Chromium at least 109.0

By: sera

Wtf is a ligature?

fi → fī

fl → fl



Wtf is a ligature?

f -> **f**

i -> **i**

fi -> 

The   ery   nch

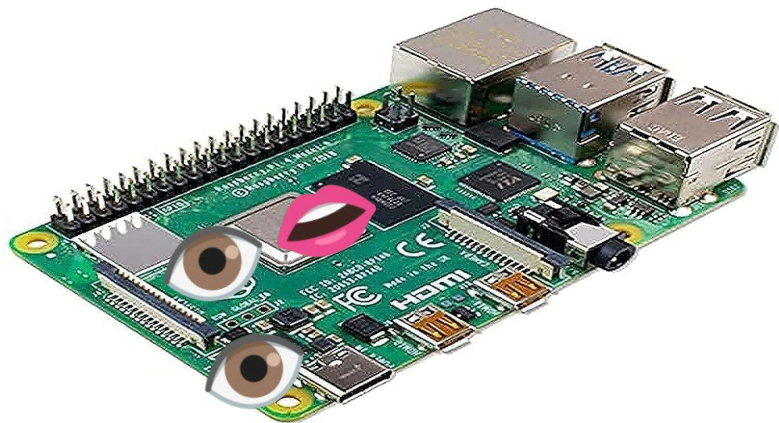
swiftly **flicked** its   ne **f**eathers

Extra Credit: Breaking Government-Approved Crypto

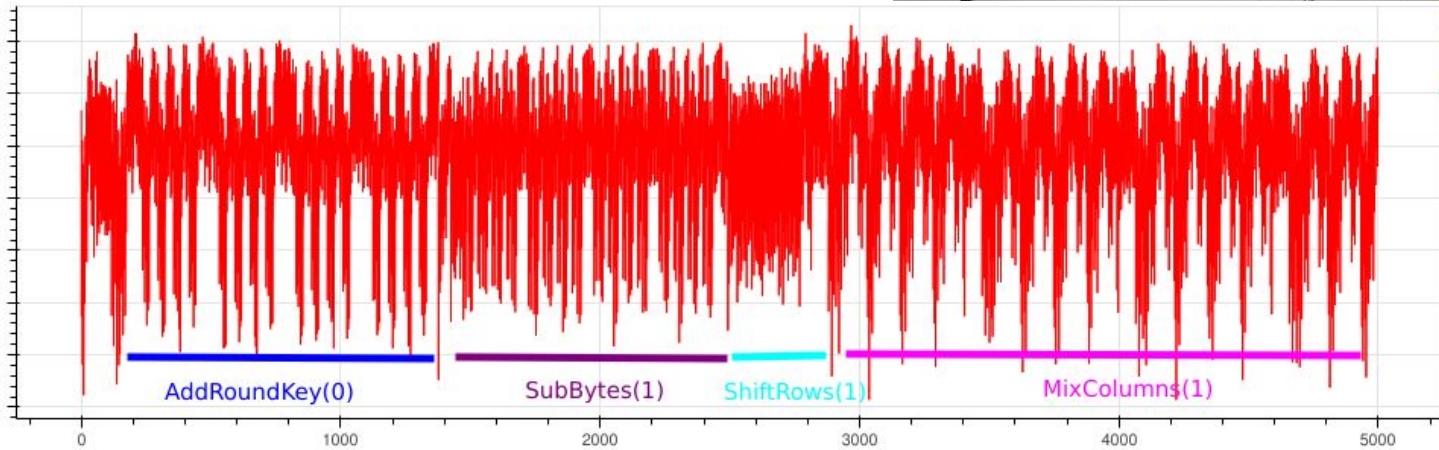
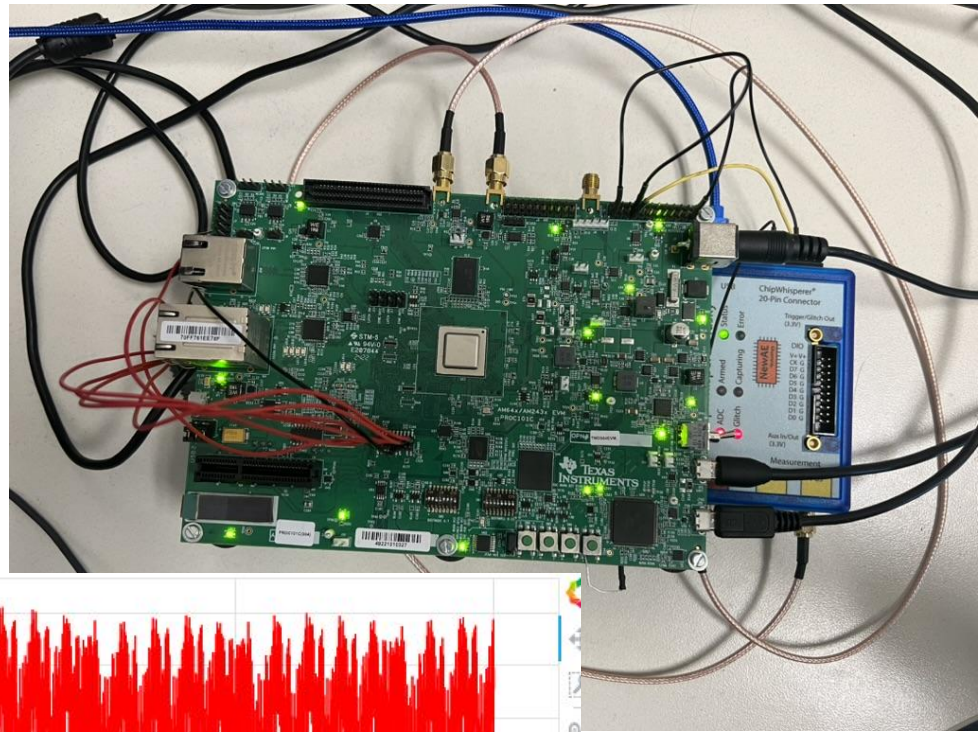
(it's brute forcing but cooler)

Breaking Government-Approved Crypto

- You have some device, it has some firmware on it
- You wanna figure out what it's doing, let's look at that firmware!
- The firmware is encrypted with AES >:(



Let's Power Analysis it!



How does it work

Simple: magic

Complex: ML magic

$$\rho_{X,Y} = \frac{\text{cov}(X,Y)}{\sigma_X \sigma_Y} = \frac{E[(X - \mu_X)(Y - \mu_Y)]}{\sqrt{E[(X - \mu_X)^2]E[(Y - \mu_Y)^2]}}$$

$$r_{i,j} = \frac{\sum_{d=1}^D [(h_{d,i} - \bar{h}_i)(t_{d,j} - \bar{t}_j)]}{\sqrt{\sum_{d=1}^D (h_{d,i} - \bar{h}_i)^2 \sum_{d=1}^D (t_{d,j} - \bar{t}_j)^2}}$$

$$r_{i,j} = \frac{D \sum_{d=1}^D h_{d,i} t_{d,j} - \sum_{d=1}^D h_{d,i} \sum_{d=1}^D t_{d,j}}{\sqrt{\left(\left(\sum_{d=1}^D h_{d,i} \right)^2 - D \sum_{d=1}^D h_{d,i}^2 \right) \left(\left(\sum_{d=1}^D t_{d,j} \right)^2 - D \sum_{d=1}^D t_{d,j}^2 \right)}}$$

Ok but for real:

- Computers take power to do stuff (duh)
- They take more power to do different operations
- $34 + 12 \rightarrow 10\text{pW}$
- $578284 + 7278894 \rightarrow 12\text{pW}$

Hamming Weight Power Model

If a number has more “1”s in it (in binary form), your CPU draws more power when using it

255 = 0b11111111 -> draws “8” power

38 = 0b00100110 -> draws “3” power

0 = 0b00000000 -> draws “0” power

The abridged version:

AES Step 1:

$\text{temp} = \text{sbox}[\text{key}[0] \wedge \text{ct}[0]]$

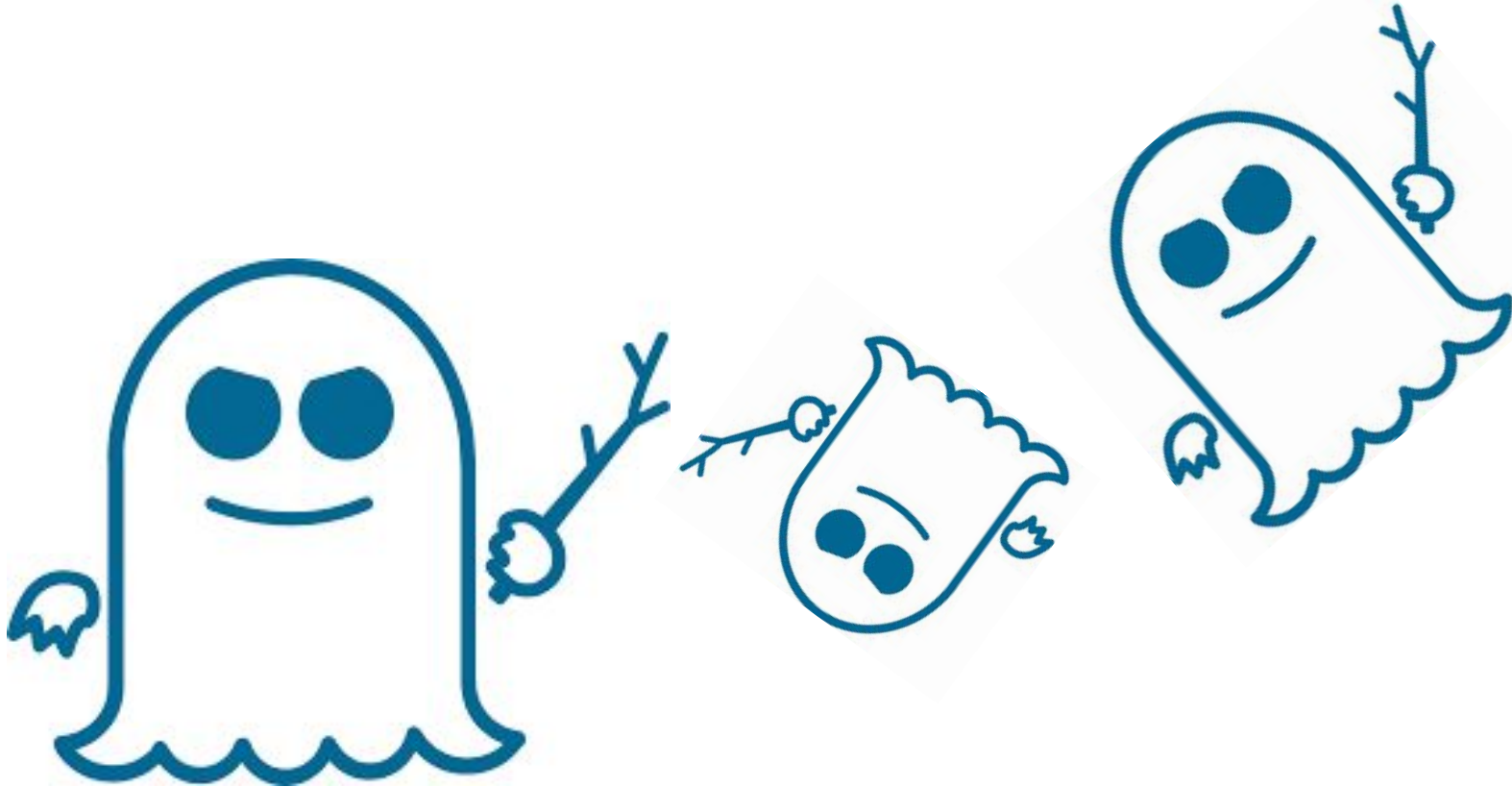
^ this is all we care
about!

	x0	x1	x2	x3	x4	x5	x6	x7	x8	x9	xa	xb	xc	xd	xe	xf
0x	63	7c	77	7b	f2	6b	6f	c5	30	01	67	2b	fe	d7	ab	76
1x	ca	82	c9	7d	fa	59	47	f0	ad	d4	a2	af	9c	a4	72	c0
2x	b7	fd	93	26	36	3f	f7	cc	34	a5	e5	f1	71	d8	31	15
3x	04	c7	23	c3	18	96	05	9a	07	12	80	e2	eb	27	b2	75
4x	09	83	2c	1a	1b	6e	5a	a0	52	3b	d6	b3	29	e3	2f	84
5x	53	d1	00	ed	20	fc	b1	5b	6a	cb	be	39	4a	4c	58	cf
6x	d0	ef	aa	fb	43	4d	33	85	45	f9	02	7f	50	3c	9f	a8
7x	51	a3	40	8f	92	9d	38	f5	bc	b6	da	21	10	ff	f3	d2
8x	cd	0c	13	ec	5f	97	44	17	c4	a7	7e	3d	64	5d	19	73
9x	60	81	4f	dc	22	2a	90	88	46	ee	b8	14	de	5e	0b	db
ax	e0	32	3a	0a	49	06	24	5c	c2	d3	ac	62	91	95	e4	79
bx	e7	c8	37	6d	8d	d5	4e	a9	6c	56	f4	ea	65	7a	ae	08
cx	ba	78	25	2e	1c	a6	b4	c6	e8	dd	74	1f	4b	bd	8b	8a
dx	70	3e	b5	66	48	03	f6	0e	61	35	57	b9	86	c1	1d	9e
ex	e1	f8	98	11	69	d9	8e	94	9b	1e	87	e9	ce	55	28	df
fx	8c	a1	89	0d	bf	e6	42	68	41	99	2d	0f	b0	54	bb	16

The abridged version:

1. Make a guess as to what the key byte is
2. Simulate what the “expected” power draw is for that key byte
3. Find how different it is from the actual power draw
4. Which guess has the closest similarity (correlation) to the actual power draw? ->
This is probably what the key byte is!
5. Repeat for next key byte

Extra Credit: Speculative Execution for dummies



Click to add title

- nah we're not getting this far not gonna make slides for this lol